Multi-Armed Bandits : an introduction

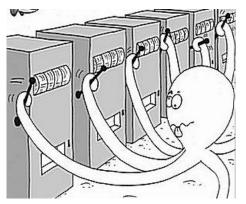
Emilie Kaufmann



EURO PhD school, July 2022

Why bandits?

one-armed bandit = old name for a slot machine



an agent facing arms in a Multi-Armed Bandit

→ How to sequentially chose which arm to pull in order to maximize our profit ? Emilie Kaufmann | CRIStAL

Sequential resource allocation

Clinical trials

▶ *K* treatment for a given symptom (with unknown effect)



Which treatment should be allocated to the next patient based on responses observed on previous patients?

Online advertisement

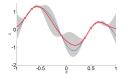
▶ K adds that can be displayed



Which add should be displayed for a user, based on the previous clicks of previous (similar) users?

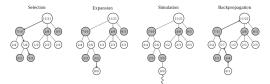
Dynamic allocation of computational resource

Numerical experiments :



where to evaluate a costly function in order to find its maximum?

Artificial intelligence for games :



how to choose the next game to simulate in order to find the best move to play next?

Outline

1 The multi-armed bandit problem

2 Fixing the greedy strategy

3 Upper Confidence Bound (UCB) algorithms

4 Towards optimal algorithms

The Multi-Armed Bandit Setting

K arms $\leftrightarrow K$ rewards streams $(X_{a,t})_{t\in\mathbb{N}}$



At round t, an agent :

- chooses an arm A_t
- ▶ receives a reward $R_t = X_{A_t,t}$

Sequential sampling strategy (bandit algorithm) :

$$A_{t+1} = F_t(A_1, R_1, \ldots, A_t, R_t).$$

Goal : Maximize $\sum_{t=1}^{T} R_t$.

The Stochastic Multi-Armed Bandit Setting

K arms \leftrightarrow *K* probability distributions : ν_a has mean μ_a



At round t, an agent :

- chooses an arm A_t
- ▶ receives a reward $R_t = X_{A_t,t} \sim \nu_{A_t}$

Sequential sampling strategy (bandit algorithm) :

$$A_{t+1}=F_t(A_1,R_1,\ldots,A_t,R_t).$$

Goal : Maximize
$$\mathbb{E}\left[\sum_{t=1}^{T} R_t\right]$$

→ a particular reinforcement learning problem

Clinical trials

Historical motivation [Thompson, 1933]



For the *t*-th patient in a clinical study,

- chooses a treatment A_t
- ▶ observes a response $R_t \in \{0,1\} : \mathbb{P}(R_t = 1 | A_t = a) = \mu_a$

Goal : maximize the expected number of patients healed

Online content optimization

Modern motivation (\$\$) [Li et al., 2010] (recommender systems, online advertisement)



For the *t*-th visitor of a website,

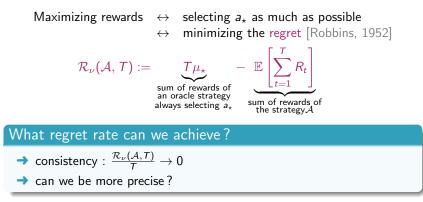
- recommend a movie A_t
- ▶ observe a rating $R_t \sim \nu_{A_t}$ (e.g. $R_t \in \{1, ..., 5\}$)

Goal : maximize the sum of ratings

Regret of a bandit algorithm

Bandit instance : $\nu = (\nu_1, \nu_2, \dots, \nu_K)$, mean of arm $a : \mu_a = \mathbb{E}_{X \sim \nu_a}[X]$.

$$\mu_{\star} = \max_{a \in \{1, \dots, K\}} \mu_a \qquad a_{\star} = \operatorname*{argmax}_{a \in \{1, \dots, K\}} \mu_a.$$



Regret decomposition

 $N_a(t)$: number of selections of arm a in the first t rounds $\Delta_a := \mu_\star - \mu_a$: sub-optimality gap of arm a

Regret decomposition

$$\mathcal{R}_{\nu}(\mathcal{A},T) = \sum_{a=1}^{K} \Delta_{a} \mathbb{E}\left[N_{a}(T)\right].$$

Proof.

$$\mathcal{R}_{\nu}(\mathcal{A}, T) = \mu_{\star} T - \mathbb{E}\left[\sum_{t=1}^{T} X_{A_{t}, t}\right] = \mu_{\star} T - \mathbb{E}\left[\sum_{t=1}^{T} \mu_{A_{t}}\right]$$
$$= \mathbb{E}\left[\sum_{t=1}^{T} (\mu_{\star} - \mu_{A_{t}})\right]$$
$$= \sum_{a=1}^{K} \underbrace{\mu_{\star} - \mu_{a}}_{\Delta_{a}} \mathbb{E}\left[\underbrace{\sum_{t=1}^{T} \mathbb{1}(A_{t} = a)}_{N_{a}(T)}\right].$$

Regret decomposition

 $N_a(t)$: number of selections of arm a in the first t rounds $\Delta_a:=\mu_\star-\mu_a$: sub-optimality gap of arm a

Regret decomposition

$$\mathcal{R}_{\nu}(\mathcal{A},T) = \sum_{a=1}^{K} \Delta_{a} \mathbb{E}\left[N_{a}(T)\right].$$

A strategy with small regret should :

- ▶ select not too often arms for which $\Delta_a > 0$
- \blacktriangleright ... which requires to try all arms to estimate the values of the Δ_a 's

\Rightarrow Exploration / Exploitation trade-off

The greedy strategy

Select each arm once and, for $t \ge K$, exploit the current knowledge :

 $A_{t+1} = \operatorname*{argmax}_{a \in [K]} \hat{\mu}_a(t)$

where

N_a(t) = ∑^t_{s=1} 1(A_s = a) is the number of selections of arm a
 μ̂_a(t) = 1/N_a(t) ∑^t_{s=1} X_s 1(A_s = a) is the empirical mean of the rewards collected from arm a

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Properties :

🖒 a simple (non-parametric) algorithm

♥ suffers linear regret

e.g. in a two armed Bernoulli bandit with means $\mu_1 > \mu_2$

$$\mathcal{R}_{
u}(\mathcal{T}) \geq (1-\mu_1)\mu_2(\mu_1-\mu_2) imes(\mathcal{T}-1)$$

Outline

- 1 The multi-armed bandit problem
- 2 Fixing the greedy strategy
- **3** Upper Confidence Bound (UCB) algorithms
- 4 Towards optimal algorithms

Given $m \in \{1, \ldots, T/K\}$,

- draw each arm m times
- compute the empirical best arm $\hat{a} = \operatorname{argmax}_{a} \hat{\mu}_{a}(Km)$
- keep playing this arm until round T

 $A_{t+1} = \hat{a}$ for $t \ge Km$

 \Rightarrow EXPLORATION followed by EXPLOITATION

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 \Rightarrow EXPLORATION followed by EXPLOITATION

Analysis for two arms. $\mu_1 > \mu_2$, $\Delta := \mu_1 - \mu_2$.

$$\begin{aligned} \mathcal{R}_{\nu}(\text{ETC},T) &= \Delta \mathbb{E}[N_2(T)] \\ &= \Delta \mathbb{E}\left[m + (T-2m)\mathbb{1}\left(\hat{a}=2\right)\right] \\ &\leq \Delta m + (\Delta T) \times \mathbb{P}\left(\hat{\mu}_{2,m} \geq \hat{\mu}_{1,m}\right) \end{aligned}$$

 $\hat{\mu}_{a,m}$: empirical mean of the first *m* observations from arm *a*

Given $m \in \{1, \ldots, T/K\}$,

- draw each arm *m* times
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- ▶ keep playing this arm until round T $A_{t+1} = \hat{a}$ for t > Km

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 $\hat{\mu}_{a,m}$: empirical mean of the first *m* observations from arm *a* \rightarrow requires a concentration inequality

Technical tool : Concentration Inequalities

Sub-Gaussian random variables : $Z - \mu$ is σ^2 -subGaussian if

$$\mathbb{E}[Z] = \mu \text{ and } \mathbb{E}\left[e^{\lambda(Z-\mu)}\right] \le e^{\frac{\lambda^2\sigma^2}{2}}.$$
 (1)

Hoeffding inequality

 Z_i i.i.d. satisfying (1). For all $s \ge 1$

$$\mathbb{P}\left(\frac{Z_1 + \dots + Z_s}{s} \ge \mu + x\right) \le e^{-\frac{sx^2}{2\sigma^2}}$$

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\Rightarrow EXPLORATION followed by EXPLOITATION

Analysis for two arms. $\mu_1 > \mu_2$, $\Delta := \mu_1 - \mu_2$.

Assumption : ν_1, ν_2 are bounded in [0, 1].

$$\mathcal{R}_{\nu}(T) = \Delta \mathbb{E}[N_2(T)]$$

= $\Delta \mathbb{E}[m + (T - 2m)\mathbb{1}(\hat{a} = 2)]$
 $\leq \Delta m + (\Delta T) \times \mathbb{P}(\hat{\mu}_{2,m} \geq \hat{\mu}_{1,m})$

 $\hat{\mu}_{a,m}$: empirical mean of the first m observations from arm a \rightarrow Hoeffding's inequality

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$$\begin{aligned} \mathcal{R}_{\nu}(T) &= \Delta \mathbb{E}[N_2(T)] \\ &= \Delta \mathbb{E}\left[m + (T-2m)\mathbb{1}\left(\hat{a}=2\right)\right] \\ &\leq \Delta m + (\Delta T) \times \exp(-m\Delta^2/2) \end{aligned}$$

 $\hat{\mu}_{a,m}$: empirical mean of the first *m* observations from arm *a* \rightarrow Hoeffding's inequality

Given $m \in \{1, \ldots, T/K\}$,

- draw each arm m times
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- keep playing this arm until round T

$$A_{t+1} = \hat{a}$$
 for $t \ge Km$

 \Rightarrow EXPLORATION followed by EXPLOITATION

$$\begin{split} & \underbrace{\text{Analysis for two arms. } \mu_1 > \mu_2, \ \Delta := \mu_1 - \mu_2. \\ & \overline{\text{Assumption : } \nu_1, \nu_2 \text{ are bounded in [0, 1].} \\ & \text{For } m = \frac{2}{\Delta^2} \log \left(\frac{T \Delta^2}{2} \right), \\ & \mathcal{R}_{\nu}(\text{ETC}, T) \leq \frac{2}{\Delta} \left[\log \left(\frac{T \Delta^2}{2} \right) + 1 \right]. \end{split}$$

Given $m \in \{1, \ldots, T/K\}$,

- draw each arm m times
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- keep playing this arm until round T

$$A_{t+1} = \hat{a} \ \ {
m for} \ \ t \geq Km$$

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Analysis for two arms. $\mu_1 > \mu_2$, $\Delta := \mu_1 - \mu_2$.

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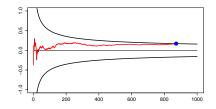
For
$$m = \frac{2}{\Delta^2} \log\left(\frac{T\Delta^2}{2}\right)$$
,
 $\mathcal{R}_{\nu}(\text{ETC}, T) \leq \frac{2}{\Delta} \left[\log\left(\frac{T\Delta^2}{2}\right) + 1\right]$

- + logarithmic regret !
- $-\,$ requires the knowledge of ${\cal T}$ and Δ

Sequential Explore-Then-Commit

explore uniformly until a random time of the form

$$au = \inf\left\{t \in \mathbb{N}: |\hat{\mu}_1(t) - \hat{\mu}_2(t)| > \sqrt{rac{c\log(\mathcal{T}/t)}{t}}
ight\}$$

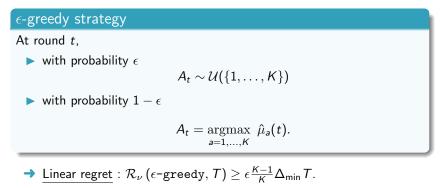


 $\blacktriangleright \ \hat{a}_{\tau} = \operatorname{argmax}_{a} \ \hat{\mu}_{a}(\tau) \text{ and } (A_{t+1} = \hat{a}_{\tau}) \text{ for } t \in \{\tau + 1, \dots, T\}$

- → [Garivier et al., 2016] for two Gaussian arms, for c = 8, same regret as ETC, without the knowledge of Δ
- → ... but larger regret as that of the best fully sequential strategy

Another possible fix : ϵ -greedy

The ϵ -greedy rule [Sutton and Barto, 1998] is a simple randomized way to alternate exploration and exploitation.



$$\Delta_{\min} = \min_{a:\mu_a < \mu_\star} \Delta_a$$

Another possible fix : e-greedy

ϵ_t -greedy strategy

At round t, • with probability $\epsilon_t := \min\left(1, \frac{K}{d^2t}\right)$ • with probability $1 - \epsilon_t$ $A_t = \underset{a=1,...,K}{\operatorname{argmax}} \hat{\mu}_a(t-1).$

Theorem [Auer et al., 2002]

$$\mathsf{lf} \ \mathsf{0} < \textit{d} \leq \Delta_{\mathsf{min}}, \ \mathcal{R}_{\nu}\left(\epsilon_t \mathsf{-greedy}, T\right) = O\left(\frac{\mathsf{K} \log(T)}{d^2}\right).$$

→ requires the knowledge of a lower bound on Δ_{\min} ...

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The optimism principle

Step 1 : construct a set of statistically plausible models

For each arm *a*, build a confidence interval on the mean μ_a :

 $\mathcal{I}_{a}(t) = [LCB_{a}(t), UCB_{a}(t)]$

LCB = Lower Confidence Bound UCB = Upper Confidence Bound

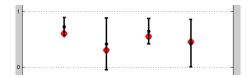


FIGURE – Confidence intervals on the means after t rounds

The optimism principle

Step 2 : act as if the best possible model were the true model (optimism in face of uncertainty)



FIGURE – Confidence intervals on the means after t rounds

▶ That is, select

$$A_{t+1} = \underset{a=1,\ldots,K}{\operatorname{argmax}} \operatorname{UCB}_{a}(t).$$

We need $UCB_a(t)$ such that

$$\mathbb{P}(\mu_a \leq \mathrm{UCB}_a(t)) \gtrsim 1 - t^{-1}.$$

→ tool : concentration inequalities

Example : rewards are σ^2 sub-Gaussian

Hoeffding inequality, reloaded

 Z_i i.i.d. satisfying (1). For all $s \ge 1$

$$\mathbb{P}\left(\frac{Z_1 + \dots + Z_s}{s} < \mu - x\right) \le e^{-\frac{sx^2}{2\sigma^2}}$$

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 Z_i i.i.d. satisfying (1). For all $s \ge 1$

$$\mathbb{P}\left(\frac{Z_1 + \dots + Z_s}{s} < \mu - x\right) \le e^{-\frac{sx^2}{2\sigma^2}}$$

Cannot be used directly in a bandit model as the number of observations from each arm is random !

N_a(t) = ∑_{s=1}^t 1_(A_s=a) number of selections of a after t rounds
 µ̂_{a,s} = ¹/_s ∑_{k=1}^s Y_{a,k} average of the first s observations from arm a
 µ̂_a(t) = µ̂_{a,Na(t)} empirical estimate of µ_a after t rounds

Hoeffding inequality + union bound

$$\mathbb{P}\left(\mu_{a} \leq \hat{\mu}_{a}(t) + \sqrt{\frac{6\sigma^{2}\log(t)}{N_{a}(t)}}\right) \geq 1 - \frac{1}{t^{2}}$$

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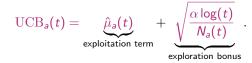
$$\mathbb{P}\left(\mu_{a} \leq \hat{\mu}_{a}(t) + \sqrt{\frac{6\sigma^{2}\log(t)}{N_{a}(t)}}\right) \geq 1 - \frac{1}{t^{2}}$$

Proof.

$$\mathbb{P}\left(\mu_{a} > \hat{\mu}_{a}(t) + \sqrt{\frac{6\sigma^{2}\log(t)}{N_{a}(t)}}\right) \leq \mathbb{P}\left(\exists s \leq t : \mu_{a} > \hat{\mu}_{a,s} + \sqrt{\frac{6\sigma^{2}\log(t)}{s}}\right)$$
$$\leq \sum_{s=1}^{t} \mathbb{P}\left(\hat{\mu}_{a,s} < \mu_{a} - \sqrt{\frac{6\sigma^{2}\log(t)}{s}}\right) \leq \sum_{s=1}^{t} \frac{1}{t^{3}} = \frac{1}{t^{2}}.$$

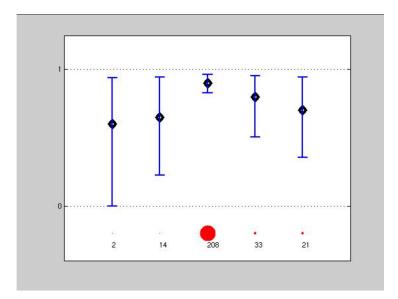
A first UCB algorithm

 $UCB(\alpha)$ selects $A_{t+1} = \operatorname{argmax}_{a} UCB_{a}(t)$ where



- popularized by [Auer et al., 2002] for bounded rewards : UCB1, for α = 2
- the analysis of UCB(α) was further refined to hold for α > 1/2 in that case [Bubeck, 2010, Cappé et al., 2013]

A UCB algorithm in action



Regret of UCB(α)

Theorem

For σ^2 -subGaussian rewards, the UCB algorithm with parameter $\alpha = 6\sigma^2$ satisfies, for any sub-optimal arm *a*,

$$\mathbb{E}_{\boldsymbol{\mu}}[N_{\boldsymbol{a}}(T)] \leq \frac{24\sigma^2}{\Delta_{\boldsymbol{a}}^2}\log(T) + 1 + \frac{\pi^2}{3}$$

where $\Delta_a = \mu_\star - \mu_a$.

Proof :



A worse-case regret bound

Corollary

$$\mathcal{R}_{\nu}(\mathrm{UCB}(6\sigma^2), T) \leq 10\sqrt{\mathcal{K}T\log(T)} + \left(1 + \frac{\pi^2}{3}\right)\left(\sum_{a=1}^{K} \Delta_a\right)$$

Proof. For any algorithm satisfying $\mathbb{E}[N_a(T)] \leq C \frac{\log(T)}{\Delta_a} + D$ for all sub-optimal arm *a*, for any $\Delta > 0$,

$$\begin{aligned} \mathcal{R}_{\nu}(T) &= \sum_{a:\Delta_{a} \leq \Delta} \Delta_{a} \mathbb{E}[N_{a}(T)] + \sum_{a:\Delta_{a} \geq \Delta} \Delta_{a} \mathbb{E}[N_{a}(T)] \\ &\leq \Delta T + \sum_{a:\Delta_{a} \geq \Delta} \left(C \frac{\log(T)}{\Delta_{a}} + D \Delta_{a} \right) \\ &\leq \Delta T + \frac{CK \log(T)}{\Delta} + D \left(\sum_{a=1}^{K} \Delta_{a} \right) \\ &= 2\sqrt{CKT \log(T)} + D \left(\sum_{a=1}^{K} \Delta_{a} \right) \text{ for } \Delta = \sqrt{\frac{CK \log(T)}{T}} \end{aligned}$$

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An improved problem-dependent result

Context : σ^2 sub-Gaussian rewards

$$UCB_{a}(t) = \hat{\mu}_{a}(t) + \sqrt{\frac{2\sigma^{2}(\log(t) + c\log\log(t))}{N_{a}(t)}}$$

(c = 0 corresponds to UCB(\alpha) with \alpha = 2\sigma^{2})

Theorem [Cappé et al.'13]

For $c \geq$ 3, the UCB algorithm associated to the above index satisfy

$$\mathbb{E}[N_a(T)] \leq \frac{2\sigma^2}{\Delta_a^2} \log(T) + C_{\mu} \sqrt{\log(T)}.$$

Summary

For UCB(α) applied to σ^2 -subGaussian reward, setting $\alpha = 2\sigma^2$ yields

▶ a problem-dependent regret bound of

$$\left(\sum_{a=1}^{K} \frac{2\sigma^2}{\Delta_a}\right) \log(T) + o(\log(T))$$

► a worse-case regret of order

$$O\left(\sqrt{KT\log(T)}\right)$$

→ how good are these regret rates?

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The Lai and Robbins lower bound

Context : a parametric bandit model where each arm is parameterized by its mean $\nu = (\nu_{\mu_1}, \dots, \nu_{\mu_K})$, $\mu_a \in \mathcal{I}$.

$$\nu \leftrightarrow \boldsymbol{\mu} = (\mu_1, \dots, \mu_K)$$

Key tool : Kullback-Leibler divergence.

Kullback-Leibler divergence

$$\mathrm{kl}(\mu,\mu'):=\mathsf{KL}\left(
u_{\mu},
u_{\mu'}
ight)=\mathbb{E}_{X\sim
u_{\mu}}\left[\lograc{d
u_{\mu}}{d
u_{\mu'}}(X)
ight]$$

Theorem

For uniformly good algorithm, $\mu_{a} < \mu_{\star} \Rightarrow \liminf_{T \to \infty} \frac{\mathbb{E}_{\mu}[N_{a}(T)]}{\log T} \ge \frac{1}{\mathrm{kl}(\mu_{a}, \mu_{\star})}$

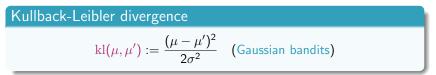
[Lai and Robbins, 1985]

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Kullback-Leibler divergence

$$\operatorname{kl}(\mu,\mu') := \mu \log\left(rac{\mu}{\mu'}
ight) + (1-\mu) \log\left(rac{1-\mu}{1-\mu'}
ight)$$
 (Be

(Bernoulli bandits)

Theorem

For *uniformly good* algorithm,

$$\mu_{a} < \mu_{\star} \Rightarrow \liminf_{T \to \infty} \frac{\mathbb{E}_{\mu}[N_{a}(T)]}{\log T} \geq \frac{1}{\mathrm{kl}(\mu_{a}, \mu_{\star})}$$

[Lai and Robbins, 1985]

UCB compared to the lower bound

Gaussian distributions with variance σ^2

- Lower bound : $\mathbb{E}[N_a(T)] \gtrsim \frac{2\sigma^2}{(\mu_\star \mu_a)^2} \log(T)$
- Upper bound : for UCB(α) with $\alpha = 2\sigma^2$

$$\mathbb{E}[N_{\mathsf{a}}(T)] \lesssim \frac{20}{(\mu_{\star} - \mu_{\mathsf{a}})^2} \log(T)$$

➔ UCB is asymptotically optimal for Gaussian rewards !

UCB compared to the lower bound

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- Lower bound : $\mathbb{E}[N_a(T)] \gtrsim \frac{2\sigma^2}{(\mu_\star \mu_a)^2} \log(T)$
- ▶ Upper bound : for UCB(α) with $\alpha = 2\sigma^2$ $\mathbb{E}[N_a(T)] \lesssim \frac{2\sigma^2}{(\mu_\star - \mu_a)^2} \log(T)$

→ UCB is asymptotically optimal for Gaussian rewards !

Bernoulli distributions (bounded, $\sigma^2 = 1/4$)

• Lower bound : $\mathbb{E}[N_a(T)] \gtrsim \frac{1}{\mathrm{kl}(\mu_a,\mu_\star)} \log(T)$

▶ Upper bound : for UCB(
$$\alpha$$
) with $\alpha = 1/2$
 $\mathbb{E}[N_a(T)] \lesssim \frac{1}{2(\mu_\star - \mu_a)^2} \log(T)$

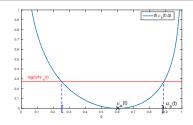
Pinsker's inequality : $kl(\mu_a, \mu_\star) > 2(\mu_\star - \mu_a)^2$

→ UCB is not asymptotically optimal for Bernoulli rewards...

The $\operatorname{kl-UCB}$ algorithm

Exploits the KL-divergence in the lower bound !

$$ext{UCB}_{a}(t) = \max\left\{q \in [0,1]: ext{kl}\left(\hat{\mu}_{a}(t),q
ight) \leq rac{\log(t)}{N_{a}(t)}
ight\}.$$



A tighter concentration inequality [Garivier and Cappé, 2011]

For Bernoulli rewards,

$$\mathbb{P}(\mathrm{UCB}_{s}(t) > \mu_{s}) \gtrsim 1 - rac{1}{t\log(t)}$$

An asymptotically optimal algorithm

kl-UCB selects $A_{t+1} = \operatorname{argmax}_{a} \operatorname{UCB}_{a}(t)$ with

$$\operatorname{UCB}_{a}(t) = \max\left\{q \in [0,1] : \operatorname{kl}\left(\hat{\mu}_{a}(t),q\right) \leq \frac{\log(t) + c\log\log(t)}{N_{a}(t)}\right\}.$$

Theorem [Cappé et al., 2013]

If $c\geq$ 3, for every arm such that $\mu_{a}<\mu_{\star}$,

$$\mathbb{E}_{\mu}[N_{a}(T)] \leq \frac{1}{\mathrm{kl}(\mu_{a},\mu_{\star})} \log(T) + C_{\mu} \sqrt{\log(T)}.$$

asymptotically optimal for Bernoulli rewards

$$\mathcal{R}_{oldsymbol{\mu}}(ext{kl-UCB}, \mathcal{T}) \simeq \left(\sum_{a: \mu_a < \mu_\star} rac{\Delta_a}{ ext{kl}(\mu_a, \mu_\star)}
ight) \log(\mathcal{T}).$$

A worse case lower bound

Theorem [Cesa-Bianchi and Lugosi, 2006]

Fix $T \in \mathbb{N}$. For every bandit algorithm \mathcal{A} , there exists a stochastic bandit model ν with rewards supported in [0, 1] such that

$$\mathcal{R}_{
u}(\mathcal{A}, T) \geq rac{1}{20}\sqrt{\kappa T}$$

worse-case model :

$$\begin{cases} \nu_a &= \mathcal{B}(1/2) \text{ for all } a \neq i \\ \nu_i &= \mathcal{B}(1/2 + \Delta) \end{cases}$$

with $\Delta \simeq \sqrt{K/T}$.

Remark. (kl)-UCB only achieves $O(\sqrt{KT \log(T)})$

Emilie Kaufmann | CRIStAL

Going further

We saw different type of frequentist algorithms :

- either based on comparing (MLE) estimates of the mean rewards (ETC, ε-greedy)
- ▶ or using confidence intervals (UCB, kl-UCB)

Next lecture : Bayesian bandits

Going further

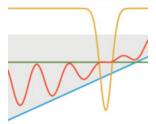
Perspectives :

- algorithms which are asymptotically optimal and minimax optimal [Garivier et al., 2018]
- algorithms which are asymptotically optimal for different families of distributions (e.g. one algorithm for Gaussian and Bernoulli bandits) [Baudry et al., 2020]
- algorithms which are robust to adversarial rewards (Best Of Both worlds)
 [Zimmert and Seldin, 2021]
- algorithms which are robust to non-stationary rewards [Garivier and Moulines, 2011, Suk and Kpotufe, 2022]

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Bandit Algorithms

TOR LATTIMORE CSABA SZEPESVÁRI



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by [Lattimore and Szepesvari, 2019]

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